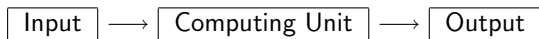


Linear Algebra and Quantum Computing

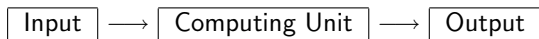
Chi-Kwong Li
Department of Mathematics
The College of William and Mary
Williamsburg, Virginia, USA

Joint work with Yiu-Tung Poon (Iowa State University).

General computing models

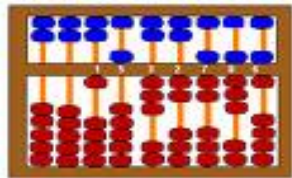


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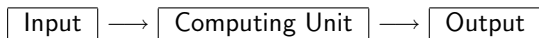


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Classical computing (Abacus)



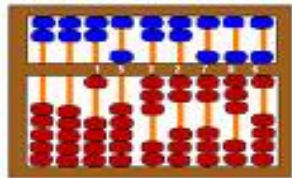
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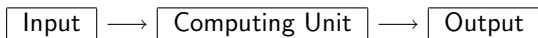
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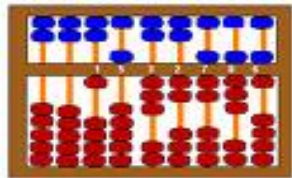
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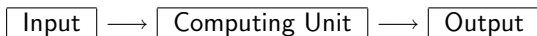
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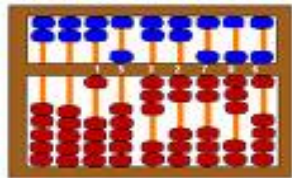
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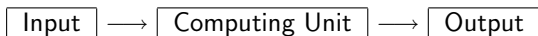
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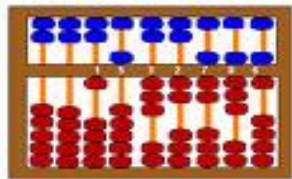
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Modern Computing (Digital Computer)



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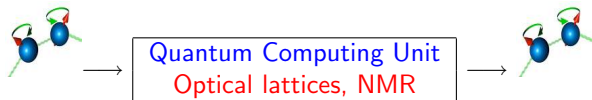
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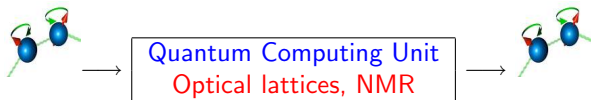
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Quantum computing

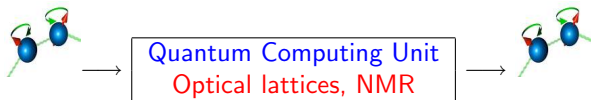


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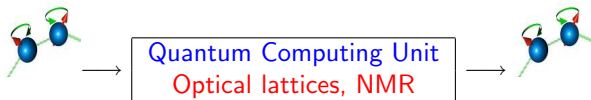
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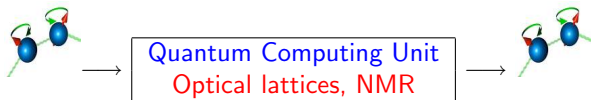
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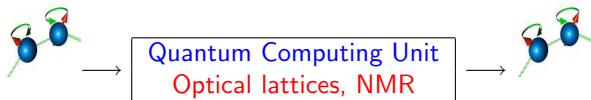
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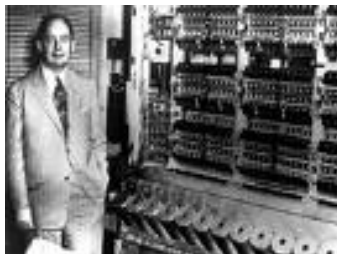
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- All these require the understanding of mathematics, physics, chemistry, computer sciences, engineering, etc.

Mathematical formulation (by von Neumann)

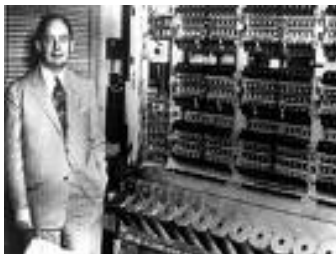
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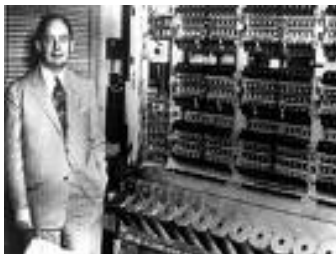
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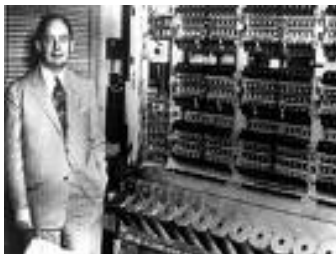
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- One can apply a quantum operation to a state in superposition.

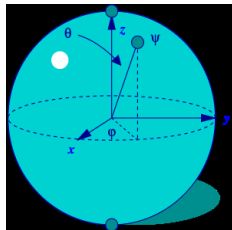


Matrix and Bloch sphere

- It is convenient to represent the quantum state $|\psi\rangle$ as a rank-one orthogonal projection:

$$|\psi\rangle\langle\psi| = \frac{1}{2} \begin{pmatrix} 1+z & x-iy \\ x+iy & 1-z \end{pmatrix}$$

with $x, y, z \in \mathbb{R}$, $x^2 + y^2 + z^2 = 1$.



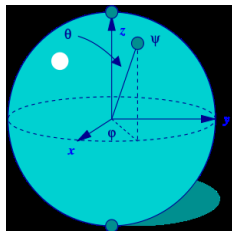
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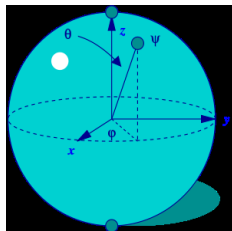
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- The theory was discovered way before the applications!



Man-Duen Choi

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A general problem

Since every quantum operation / channel is a trace preserving completely positive linear map, it is interesting to study the following.

Question

Given $A_1, \dots, A_k \in M_n$ and $B_1, \dots, B_k \in M_m$, is there a (unital/trace preserving) completely positive linear map \mathcal{L} satisfying

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- Understand the **duality** relation between the trace preserving completely positive linear maps and the unital preserving completely positive linear maps.

Basic results

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It is known that $x \prec y$ if and only if there is a **doubly stochastic** matrix D such that $x = yD$.

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- One can use D to construct $m \times n$ matrices F_1, \dots, F_r with $r = \max(m, n)$ such that

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- $\lambda(B) \prec (a_+, 0, \dots, 0, a_-)$ in $\mathbb{R}^{1 \times m}$.
- There is an $n \times m$ **row stochastic matrix** (nonnegative matrix with all row sums equal to one) D such that $\lambda(B) = \lambda(A)D$.

- The matrix D can be chosen so that the first k rows all equal and the last $n - k$ rows all equal.
- One can use D to construct $m \times n$ matrices F_1, \dots, F_r with $r = \max(m, n)$ such that

$$B = \sum_{j=1}^r F_j A F_j^* \quad \text{and} \quad \sum_{j=1}^r F_j^* F_j = I_n.$$

- For density matrices A and B , the condition trivially holds.

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Question

Can we deduce this result from the previous one using **duality** of completely positive linear map?

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Suppose $A = \text{diag}(4, 1, 1, 0)$ and $B = \text{diag}(3, 3, 0, 0)$. Then there is a trace preserving completely positive map sending A to B ,

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Reason: If there were such a map, it would send A_1 to B_1 for

$$A_1 = A - I_4 = \text{diag}(3, 0, 0, -1) \quad \text{and} \quad B_1 = B - I_4 = \text{diag}(2, 2, -1, -1),$$

which is a contradiction.

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- B is in the convex hull of the **unitary orbit** $\mathcal{U}(A)$ of A :

$$\{UAU^* : U \text{ unitary}\}.$$

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- The problem reduces to **joint/multivariate majorization**.
- What about non-commuting families?

CP maps with restricted Kraus (Choi) rank

For given $A \in H_n$, $B \in H_m$, and a positive integer r , we are interested in **constructing/finding** a CP map $\mathcal{L} : M_n \rightarrow M_m$ of the form $\mathcal{L}(A) = \sum_{j=1}^r F_j A F_j^*$ such that $\mathcal{L}(A) = B$.

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- The number of positive (negative) eigenvalues of $A \otimes I_r$ is more than or equal to the number of positive (negative) eigenvalues of B .

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Remark Even if $A, B \in H_n$ are density matrices, the roles of A and B are not symmetric in the last two theorems.

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Remark Conditions (b) and (c) can be expressed in terms of eigenvalue inequalities using **Littlewood-Richardson rules** or the **Horn sequences**:

$$\sum_{j \in J_0} \lambda_j(B) \leq \sum_{j \in J_1} \lambda_j(C_1) + \dots + \sum_{j \in J_r} \lambda_j(C_r)$$

with $(J_0, J_1, \dots, J_r) \in \mathcal{S}(n, r, \ell)$ for $\ell \in \{1, \dots, n-1\}$.

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Remark This is not true for unital completely positive linear maps.

A computational approach

Derive **numerical scheme** (using gradient flow, positive semi-definite programming, etc.) to solve the following:

Given $A_1, \dots, A_k \in H_n$, $B_1, \dots, B_k \in H_m$, determine \mathcal{L} such that

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$$\mathcal{L}(A_j) = B_j, \quad j = 1, \dots, k,$$

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- The sum of $r \times r$ principal submatrix of \mathcal{L} : $S_r(\mathcal{L}) = 0$ for a given r (\mathcal{L} has rank less than r).

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Thank you for your attention!